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- A very brief description of short-read alignment
- Arioc: a GPU-accelerated short-read aligner
- What is a "large" genome?
- A software view of a reference genome
- Repetitiveness versus speed
- Performance

Short-read alignment

Perfect alignment

- R: CATGTGTGAAGCCTCCATACTTGAGTCCTGAACTGATGAACTAA
- Q:

AAGCCTCCATACTTGAGTCCTGAACTGATGAA

Alignment with mismatches

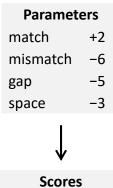
R: CATGTGTGAAGCCTCCATACCTGAGTCATGAACTGATGAACTAA

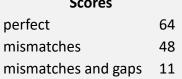
Alignment with mismatches and gaps

R: CATGTGTGAAGCCGCGCGCCCATACATGAGTCATGAAC--ATGAACTAA

Q: AAGCCT----CCATACTTGAGTCCTGAACTGATGAA

Scoring example





Short-read alignment

Extract and hash subsequences ("seeds")

Q: AAGCCTCCATACTTGAGTCCTGAACTGATGAA AAGCCTCCAT → 0xDEA5D502

AGULTULAT	7	0XDEA5D502
AGCCTCCATA	\rightarrow	0x29DEC1F0
GCCTCCATAC	\rightarrow	0xDB840577
CCTCCATACT	\rightarrow	0x4DBA90D5
• • •		

Probe hash table to find reference-sequence locations

0xDEA5D502: 01:14353363, 01:15536663, 02:06335366 ... 0x29DEC1F0: 01:14353364, 06:20159342, 18:00513566 0xDB840577: 01:14353365, 01:15536665, 05:83754151 ... 0x4DBA90D5: (none)

Look for high-scoring alignments ("extend") at high-priority reference-sequence locations

R: CATGTGTGAAGCCGCCATACCTGAGTCATGAAC--ATGAACTAA

Q: AAGCCTCCATACTTGAGTCCTGAACTGATGAA

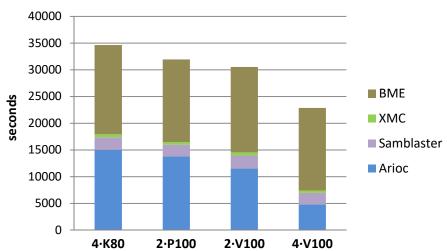
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Arioc: a GPU-accelerated short-read aligner

- Speed
 - Short-read alignment is just one step in a processing "pipeline"; the idea is that this step should not be a bottleneck
 - Order-of-magnitude (~10x) faster than CPU-only implementations
- Sensitivity
- Accuracy
- Capable of handling real-world data
 - Full-sized sequencer runs
 - Human reference genome (and larger)

Arioc is fast

- 1,304 WGBS samples
 - 150bp paired-end
 - Human reference genome
 - Average sample size: 487,757,780 pairs (975,515,560 reads)
- One step in a series of analysis tools
 - Arioc
 - Samblaster
 - Bismark methlylation extractor
- Shared compute nodes at MARCC (Maryland Advanced Research Computing Center)



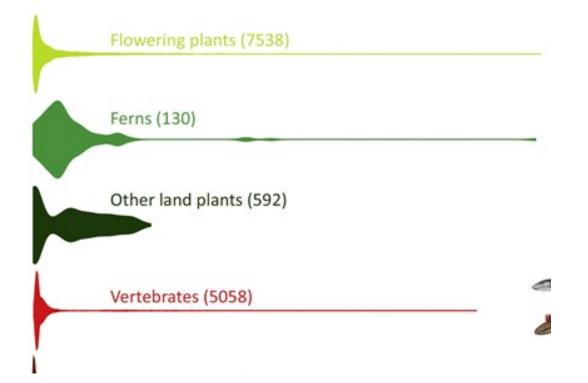
Average elapsed time per sample

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"Large" compared to what?

- The human genome is a good starting point for comparison
 - About 3 billion nucleotide bases
 - If you number each base position consecutively, you can identify each base with a 32-bit integer!
- Some interesting organisms have genomes that contain much more DNA than does the human genome

What is a large genome?



Some large genomes whose DNA has been sequenced

Organism	Size (×10 ⁹)
Mexican axolotl	32
Pine tree	22
Wheat	14.5
Human	3.2
Mouse	2.7

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Identifying genome locations

Chromosomes in axolotl genome

0.66

14

•	Subunit ID	Chromosome	Size (×10 ⁹)
	 Usually a chromosome number 	1Q	1.48
		2P	1.41
	Range of values: 1-127	2Q	1.51
		3P	1.24
	DNA strand	3Q	1.26
	DINA Sulatiu	4P	1.16
	 Forward or roverse complement 	7	2.03
	 Forward or reverse complement 	4Q	1.29
	Range of values: 0-1	8	1.71
	- Range Of Values. 0-1	5P 9	1.29
		5Q	1.50
	Offset from the start of the DNA sequence	10	1.64
		6P	1.55
	Range of values: 0-2,147,483,647	11	1.44
		6Q	1.59
		12	1.21
		13	0.72

Reference genome position in C++

```
/* 40-bit (5-byte) representation of a J value */
struct Jvalue5
{
    enum bfSize
    {
        bfSize J =
                       31, // 0..30: J (0-based offset into reference sequence)
        bfSize s =
                        1, // 31..31: strand (0: R+; 1: R-)
        bfSize subId = 7, // 32..38: subId (e.g., chromosome number)
        bfSize_x =
                        1 // 39..39: end-of-list flag
    };
    enum bfMaxVal : UINT64
    {
        bfMaxVal_J =
                         (static_cast<UINT64>(1) << bfSize_J) - 1,</pre>
        bfMaxVal s =
                         (static_cast<UINT64>(1) << bfSize_s) - 1,</pre>
        bfMaxVal_subId = (static_cast<UINT64>(1) << bfSize_subId) - 1,</pre>
        bfMaxVal x =
                         (static_cast<UINT64>(1) << bfSize_x) - 1</pre>
    };
    UINT32 J
                  : bfSize J;
    UINT32 s
                  : bfSize_s;
    UINT8 subId : bfSize subId;
    UINT8 X
                  : bfSize_x;
};
```

Large genome \rightarrow large lookup tables

Extract and hash subsequences ("seeds")

Q: AAGCCTCCATACTTGAGTCCTGAACTGATGAA

AAGCCICCAI	\rightarrow	0xDEA5D502
AGCCTCCATA	\rightarrow	0x29DEC1F0
GCCTCCATAC	\rightarrow	0xDB840577
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Look for high-scoring alignments ("extend") at high-priority reference-sequence locations

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Q: AAGCCTCCATACTTGAGTCCTGAACTGATGAA

Hash table data-sort sizes

32-bit seeds	# lists to sort	# locations to sort
human	1,263,683,062	3,687,638,902
wheat	2,120,243,009	20,602,998,718

"Sortable" reference genome position in C++

```
/* 64-bit (8-byte) representation of a 40-bit (5-byte) J value */
struct Jvalue8
{
    enum bfSize
    {
        bfSize_J =
                       31,
                               // 0..30: J (0-based offset into reference sequence)
        bfSize s =
                        1,
                               // 31..31: strand (0: R+; 1: R-)
        bfSize_subId = 7, // 32..38: subId (e.g., chromosome number)
                      1, // 39..39: flag (used only for sorting and filtering J lists; zero in final J table)
        bfSize x =
                               // 40..63: used for sorting (see tuSortJgpu)
        bfSize_tag = 24
    };
    enum bfMaxVal : UINT64
    {
        bfMaxVal_J =
                         (static_cast<UINT64>(1) << bfSize_J) - 1,</pre>
        bfMaxVal s =
                         (static cast<UINT64>(1) << bfSize s) - 1,</pre>
        bfMaxVal subId = (static cast<UINT64>(1) << bfSize subId) - 1,</pre>
        bfMaxVal x =
                         (static_cast<UINT64>(1) << bfSize_x) - 1,</pre>
        bfMaxVal_tag = (static_cast<UINT64>(1) << bfSize_tag) - 1</pre>
    };
    UINT64 J
                  : bfSize_J;
    UINT64 s
                  : bfSize_s;
    UINT64 subId : bfSize subId;
    UINT64 x
                  : bfSize_x;
    UINT64 tag : bfSize tag;
};
```

A bit-packed segmented sort

The lists are sorted in a call to a CUDA Thrust sort implementation

/* Sort the current J-list buffer chunk. Since each 64-bit value contains a "tag" that associates the value with the J list that corresponds to an H (hash key) value, this is in effect a segmented operation. */ thrust::device_ptr<UINT64> ttpJbuf(m_pJbuf->p); thrust::sort(epCGA, ttpJbuf, ttpJbuf+m_pJbuf->Count);

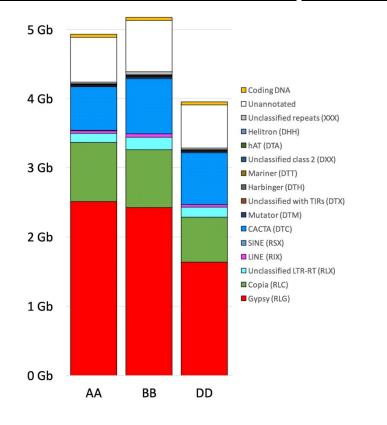
- The high-order bits identify individual lists so the result is effectively a segmented sort
- There are more lists than can be uniquely identified in the available highorder bits, so the Thrust sort API is called iteratively

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Large genomes are repetitive

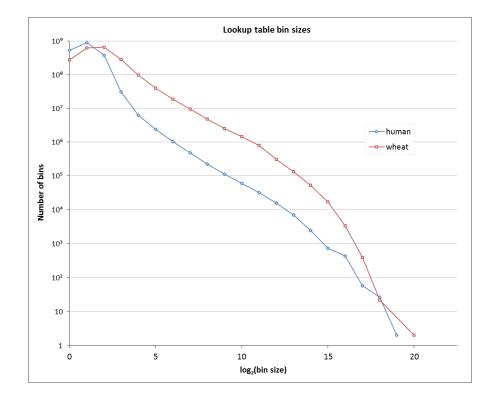
- The genome of complex organisms contains mostly non-coding DNA
- Non-coding DNA in large genomes contains many long and/or repetitive sequences
- The repetitive nature of the DNA affects the way we construct the hash tables (lookup tables)

Large genomes are repetitive



- In large genomes, much of the DNA is repetitive:
 - Human: ~50%
 - Bread wheat (Triticum aestivum): ~85%
- Repetitive DNA may contain...
 - Multiple copies of a variety of short subsequences
 - Low-information sequences
 - Homopolymers (e.g., AAAAAAAAAAAAA)
 - Tandem repeats (e.g., CGCGCGCGCGCG)
- With a large repetitive genome, any given read may have multiple locations at which it aligns
 - More alignment computations per read
 - Increased post-alignment processing to identify and classify high-scoring mappings for each read

Repetitive genome \rightarrow large lookup tables

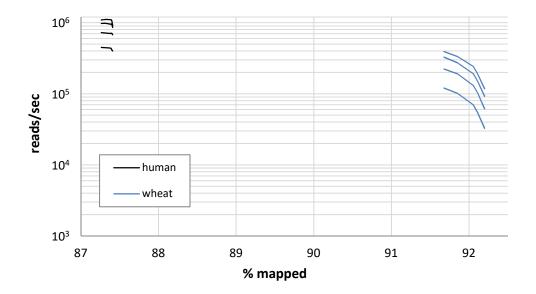


- Large-genome lookup tables (LUTs) contain more reference-sequence locations per hash value
 - Big LUT bins → more alignments computed
- Large-genome LUTs are hard to optimize
 - Pruning highly-repetitive seed locations decreases sensitivity in the read aligner

nJ	human	wheat	
raw	5,875,619,304	28,516,821,874	
final nJ	5,261,735,533	28,516,821,874	
% pruned	10%	0%	

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Large genome \rightarrow more work



human	ERR1347712	Simons Foundation Genome Diversity Project: SA_Kusunda_K-15_M
wheat	SRR6001710	Sequencing of flow sorted chromosome 7D from Canthatch K

- ~5x slower
 - Wheat vs human
 - Speed vs sensitivity
 - One through four V100 GPUs

Large genome \rightarrow more work

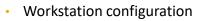
- More data transferred to GPU
 - T_R/T_Q = (time to load reference sequences)
 ÷ (time to load reads)
- More alignment computations per read
 - tuAlignN* = nongapped aligner (spaced seeds)
 - tuAlignG* = gapped aligner (Smith-Waterman)

	human	wheat
T _R /T _Q	1.22	13.32

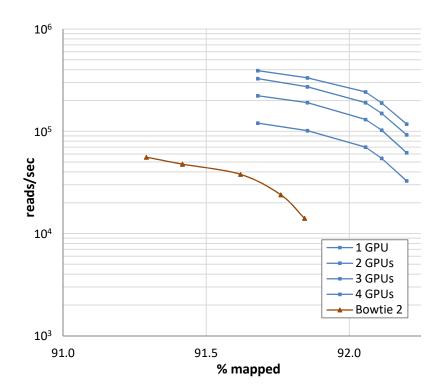
	human	wheat
tuAlignN20	3.56	17.66
tuAlignN52	10.31	113.97
tuAlignGs22	1.86	31.79
tuAlignGs12	0.65	23.67
tuAlignGwn12	0.36	2.60
tuAlignGs42	0.03	0.25

Speed versus sensitivity

- GPU-accelerated implementation is about 10x faster than CPU-only implementation
 - Arioc (2019)
 - Bowtie 2 (2019)
- GPU-accelerated implementation scales appropriately with multiple GPUs

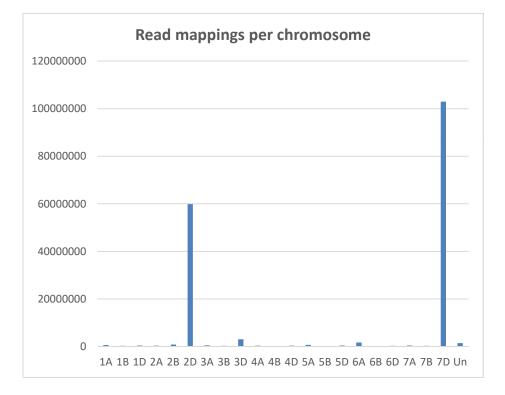


- 4 × Intel Xeon Gold 6130 CPU @ 2.10GHz (64 total threads)
- 384GB RAM
- CentOS 7.5.1804
- 4 × Nvidia V100, 32GB RAM each



Accuracy

- WGS sample SRR6001710 contains DNA from a subset of the wheat genome (flow-sorted chromosome 7D)
- Distribution of mapped reads is reasonable
 - Chromosome 2D is a major contaminant (its size is very close to that of chromosome 7D)
 - Very similar to Bowtie 2



Takeaways

- A short-read aligner runs slower with a large reference genome, but not prohibitively so
- The size of the genome demands capable hardware:
 - Up to 500GB of system RAM (for building lookup tables)
 - Fast GPUs (Volta microarchitecture at a minimum)
- The repetitive nature of the genome requires careful software configuration:
 - Do not "over-optimize" the lookup tables
 - Place some of the lookup tables in GPU RAM
 - Choose optimal criteria for concordant (proper) mappings and optimal speed-vssensitivity tradeoff



Arioc is available on Github: https://github.com/RWilton/Arioc

